# **Analyzing Access Control Overrides**

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## Aniketos in a Nutshell

Aniketos: Make composite services able to establish and maintain security and trustworthiness

## **Goals of the Aniketos platform:**

- Design-time discovery, composition and evaluation, threat awareness
- Runtime adaptation or change in service configuration
- Runtime monitoring, detection, notification

### Two related dimensions:

- Trustworthiness: Reputation, perception, centralized vs. distributed
- Security properties: Behavior, contracts, interfaces, formal verification



http://www.aniketos.eu

#### **Aniketos Fact-Sheet:**

- EU Integrated Project (IP), FP7 Call 5
- Budget: € 13.9 Mio (€ 9.6 Mio funding)
- 42 month (Aug. 2010 Feb. 2014)
- Coordinator: Sintef (Norway)
- Consortium: ATC, ATOS, DAEM, DBL, ELSAG, CNR, ITALTEL, LJMU, SAP, SEARCH-LAB, TECNALIA, THALES, TSSG, UNITN, PLUS, WIND

**SAP** applies and develops *formal* methods for ensuring the security and technical trustworthiness of services.





## **Outline**

1 Motivation

2 An Architecture Supporting for Analyzing Access Control

3 Analyzing Access Control Overrides

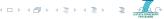
4 Conclusion

## **Our Vision**

### Assume,

- we are a nurse
- trying to access the patient record of Peter Meier . . .





## **Our Vision**







## **Our Vision**







### Motivation

### **Overriding Access Control**

- Allows the user to temporarily extend his permissions
- Also known as Break-Glass or Break the Glass (BTG)

### Relies on a post-hoc audit to evaluate the override

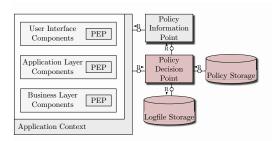
- Effort for auditing overrides increases costs
- Support auditor to reduce time and effort





## Standard Architecture

- Multiple PEPs accessing a central PDP
- Policies are loaded from a Policy Storage
- Policy Information Point (PIP) to resolve information from the application context
- Access control requests and results are stored in a Logfile Storage



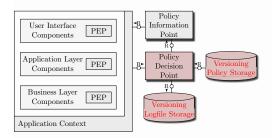




## Versioning

### Based on XACML

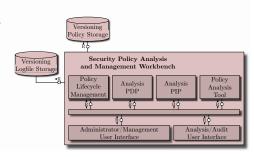
- Store policies in a Versioning Policy Storage
- Save all PIP-resolved data in a Versioning Logfile Storage
  - XACML: resolved attributes
  - Save the current "state" of the system as seen by the PDP
- Interface for clients and PIP remains the same





# **Analysis Workbench**

- Analysis PDPs load any policy version from the policy store
- Analysis Policy Information Point (PIP) as context provider
  - Analysis PIP retrieves attributes from log store
  - Simulated runtime environment for analysis
- Replay (re-evaluate) recorded (or new) access control requests
- XACML engine analysis enhancements allow for advanced analysis methods
  - Debugging of Policies
  - Abstract evaluation
  - Policy animation

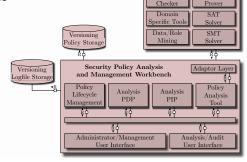






# External Analysis Tools

- Integrate existing and new developed tools
- Provide interfaces to access policy and log store
- Load and use Analysis PDPs
  - Define or modify the simulated runtime environment
  - Retrieve evaluation events from the Analysis PDP
  - Browse the evaluation state







External Analysis Tools

# Replay Access Control Requests

### To replay an access control request

- Select log entry from the log store
- Instantiate an Analysis PDP with a policy version
- Replay request on Analysis PDP
- Analysis PDP retrieves attributes as recorded for this request via Analysis PIP from the log store

### Support for understanding policies changes, e.g.,

- Replaying incidents or suspicious requests with different policy versions
  - Does a change in the policy lead to a different result?





# NHS policies in XACML

# Security policy for a NHS (National Health Service) electronic health record service (Becker, 2005):

- Permissions rely on relationships
- Policy how a relationship can be established

### Modeling in XACML

- Relationships as attributes, e.g.,
  - Patient has a set of treating clinicians
- Saved as part of the policy in the policy store
- Resolved at runtime by an XACML attribute designator





# **Policy Administration**

- Policy for management of relationships can be seen as administration policy
- Application to manage relationships can be seen as Policy Administration Point (PAP)
- PAP application has to enforce the administration policy, e.g.,
  - Who is permitted to add relationships
  - Implement obligations, e. g.,"on delete cascade" for relationships (e. g., referral)



# Versioning of policies

### Versioning of XACML policies

- Subversion (svn) as versioning system for XML files
- Logging active policy version

### Versioning of attributes: save to database

- Validity (i. e., from to)
- Depending entity (or entities), e.g.,
  - Treating clinicians depend on patient id
- Type and further information of relationship, e.g.,
  - Patient assigning a treating clinician,
    - Patient can revoke relationship
  - Treating clinician referring patient and assigning referred clinical
    - Patient cannot revoke referral relationship
    - Save referring clinical





### **Break-Glass Scenario**

In an emergency situation, there may be no valid patient - clinician relationship, e.g.,

- Patient is unconscious or not able to confirm a relationship
- No agent (i. e., trusted person of the patient) is available at time to confirm a relationship
- Due to an overwhelming emergency situation, a required referral is not entered to the IT system instantly

But, a clinician requires access to the patients health record

- Clinician uses Break-Glass to access the required data
- The access is marked as emergency access and has to be evaluated in a post-hoc audit phase





# Post-hoc Analysis I

Analyzing Simple Access Control Overrides

### After the emergency situation

- The patient is (hopefully) able to confirm the relationship
- The referral is entered to the system

### Using our replay approach, an auditor can easily

- Load a PDP with a policy version from, e.g., twenty four hours later
- Replay the accesses in questions against this policy version
- Information not available at access time can be used to verify the Break-Glass access post-hoc in a semi-automated fashion



# Post-hoc Analysis II

Process-based Compliance Checks

### **Observation:**

Many compliance regulations cannot be directly mapped to access control policies.

### **Problem:**

After overriding a single access control decision, it is unclear which compliance goals might be violated.

### Idea:

- Use (process) mining techniques for re-constructing the actual
  - process executed and
  - data-flow

that took place.

- Apply formal analysis techniques (e.g., using AVANTSSAR tools)
  - determine the set of high-level compliance and security requirements that were violated by the overridden access control decision.



### Conclusion

### Break-Glass allows to

Write restrictive policies for the regular case, as, in emergency situations, access is still possible

#### Our framework allows to

- Use information at post-hoc time which was not available at access time
- Enable (partially) semi-automated evaluation of Break-Glass accesses and therefore reduce effort and costs





# Thank you for your attention!

Any questions or remaks?

## **Related Publications**



Achim D. Brucker and Helmut Petritsch.

Extending access control models with break-glass.

In Barbara Carminati and James Joshi, editors, ACM SACMAT, pages 197–206. acm Press, 2009.

ISBN 978-1-60558-537-6.

http://www.brucker.ch/bibliography/abstract/brucker.ea-extending-2009.



Achim D. Brucker and Helmut Petritsch.

Idea: Efficient evaluation of access control constraints.

In F. Massacci, D. Wallach, and N. Zannone, editors, *International Symposium on Engineering Secure Software and Systems (ESSoS)*, number 5965 in Incs, pages 157–165. Springer, 2010.

ISBN 978-3-642-11746-6.

http://www.brucker.ch/bibliography/abstract/brucker.ea-efficient-2010.



Achim D. Brucker and Helmut Petritsch.

A framework for managing and analyzing changes of security policies.

In IEEE POLICY, pages 105–112. ieee Computer Society, June 2011.

ISBN 978-0-7695-4330-7/11.

http://www.brucker.ch/bibliography/abstract/brucker.ea-framework-2011.



PoFI 2011